

Completeness of Iris-based Program Logics



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with

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A Missing Piece in Iris

Logic

Truth

A Missing Piece in Iris

Logic

$\text{safe}_\varphi(e, \sigma)$

Truth

Never gets stuck
Result satisfies φ

A Missing Piece in Iris

$\vdash \text{wp } e \{v. \ulcorner \varphi(v) \urcorner\}$

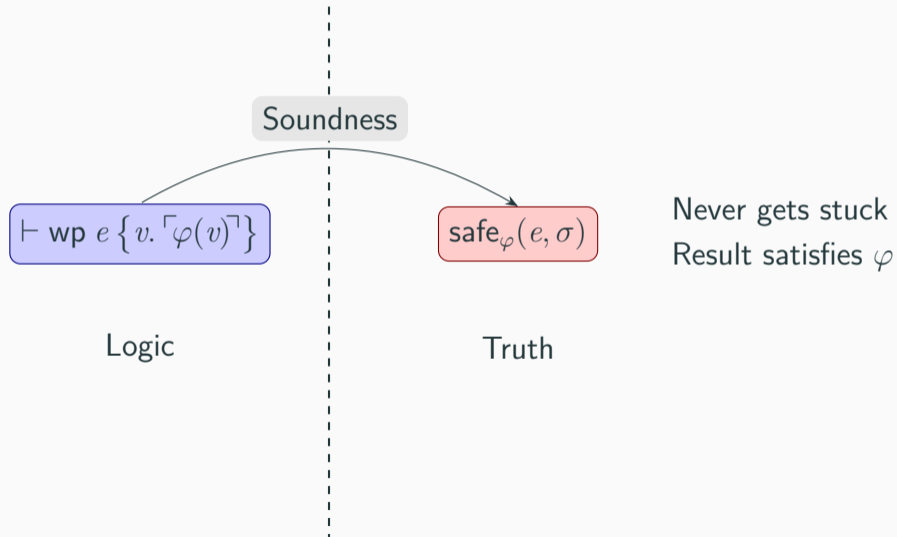
Logic

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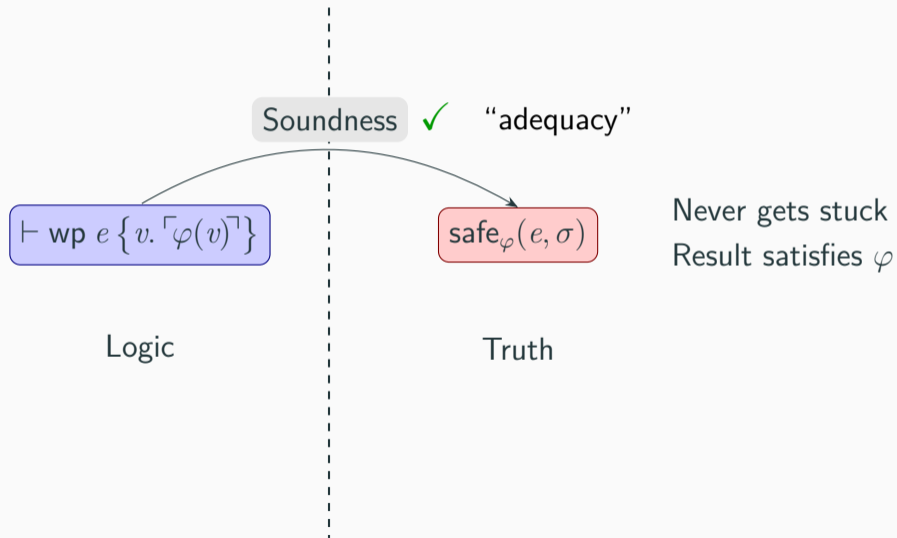
Truth

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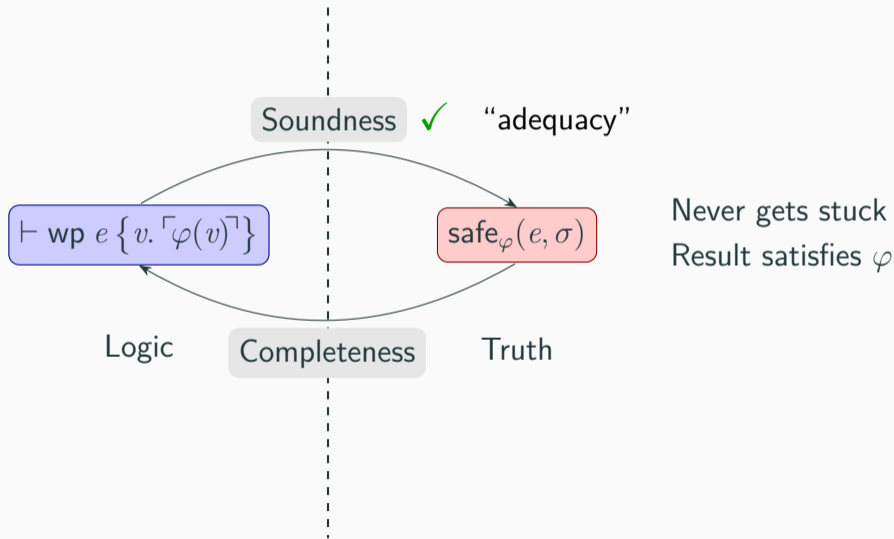
A Missing Piece in Iris



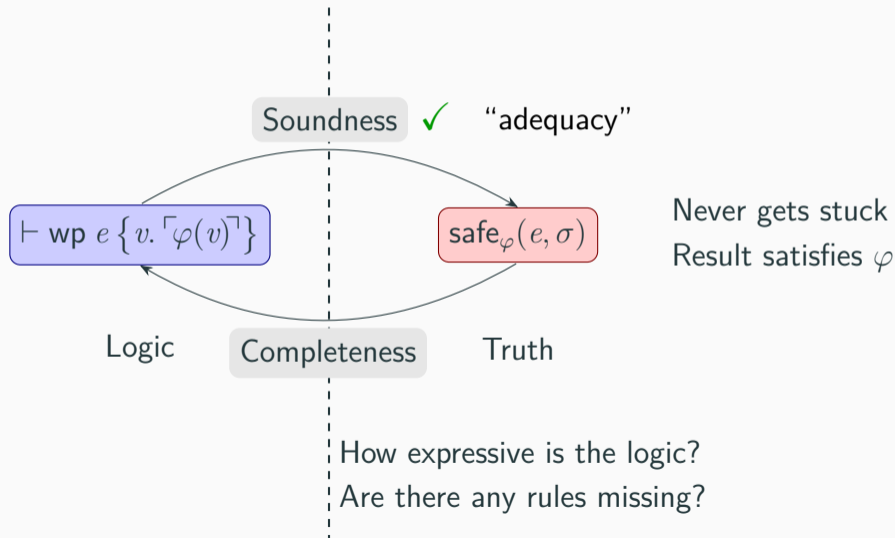
A Missing Piece in Iris



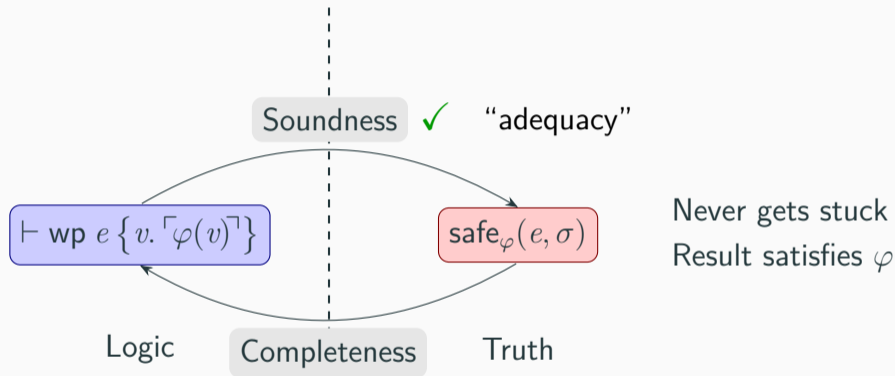
A Missing Piece in Iris



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A Missing Piece in Iris

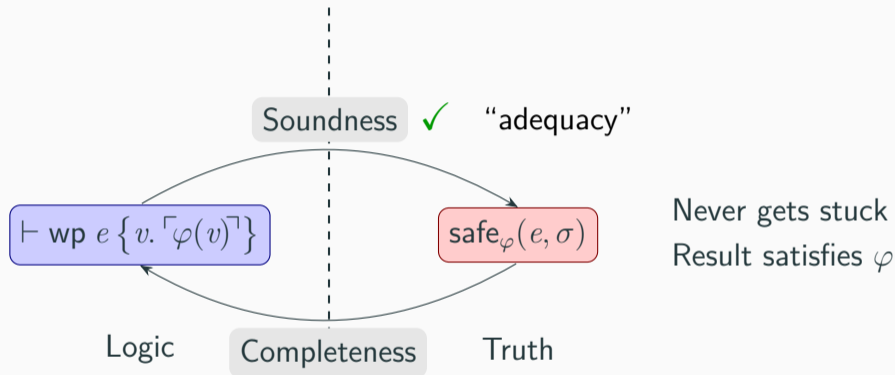


Prior work:

Cook (1978), Owicki (1975),
Jones (1983), Stirling (1988), etc.

How expressive is the logic?
Are there any rules missing?

A Missing Piece in Iris



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Iris: ?

How expressive is the logic?
Are there any rules missing?

Our Contributions

Part 1: completeness of whole program specifications

- Generic completeness theorems for any Iris-based program logics for
 - Partial correctness
 - Total correctness
 - Refinement
- Concrete languages:
 - HeapLang, λ_{Rust} , Time credits, Eris
 - **Probably your favorite program logics!**

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- Generic completeness theorems for any Iris-based program logics for
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 - **Probably your favorite program logics!**

Part 2 (WIP): completeness of library specifications

- Completeness of logically atomic triples in Iris

Whole Program Completeness

Johannes Hostert

Zichen Zhang

Puming Liu

Simon Gregersen

Ralf Jung

Joseph Tassarotti

Warm-up: Sequential Completeness

Language: $\rightarrow_{\text{pure}} \mid \mathbf{ref} \ e \mid !e \mid e_1 \leftarrow e_2$

HeapLang, but without **fork**, **free**, and prophecy variables

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$$\forall e \ . \text{safe}_\varphi(e, \emptyset) \implies \vdash \text{wp} \ e \ \{v. \lceil \varphi(v) \rceil\}$$

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$$\forall e, \sigma. \text{safe}_\varphi(e, \sigma) \implies S_o(\sigma) \vdash \text{wp } e \{v. \ulcorner \varphi(v) \urcorner\}$$

$$S_o(\sigma) \triangleq \bigstar_{(l \leftarrow v) \in \sigma} (l \mapsto v) \quad (\text{the ownership of } \sigma)$$

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Löb induction $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap S_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

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- If $\text{safe}_\varphi(v, \sigma)$ ($v \in \text{Val}$)
- If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$),

Warm-up: Sequential Completeness

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- If $\text{safe}_\varphi(v, \sigma)$ ($v \in \text{Val}$) $\implies \lceil \varphi(v) \rceil \vdash \text{wp } v \{v. \lceil \varphi(v) \rceil\}$ (by WPVALUE)
- If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$),

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- If $\text{safe}_\varphi(v, \sigma)$ ($v \in \text{Val}$) $\implies \lceil \varphi(v) \rceil \vdash \text{wp} \ v \ \{v. \lceil \varphi(v) \rceil\}$ (by WPVALUE)
- If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$), then e_1 is reducible

case analysis on the first step

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp $\text{Hsafe1: safe}_\varphi(e_1, \sigma_1)$

Persistent hyp $\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp $\text{Ho: S}_o(\sigma_1)$

Goal $\text{wp } e_1 \{v. \lceil \varphi(v) \rceil\}$

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

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Goal $\text{wp } e_1 \{v. \lceil \varphi(v) \rceil\}$

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp $\text{Hsafe1: safe}_\varphi(e_1, \sigma_1)$

$\text{Hprim: } (e_1, \sigma_1) \rightarrow_{\text{prim}} (e_2, \sigma_2)$

$\text{Hsafe2: safe}_\varphi(e_2, \sigma_2)$

Persistent hyp $\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

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If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp	$\text{Hsafe1: safe}_\varphi(e_1, \sigma_1)$	unfold prim-step	$e_1 = K[e'_1] \wedge e_2 = K[e'_2]$
	$\text{Hprim: } (e_1, \sigma_1) \rightarrow_{\text{prim}} (e_2, \sigma_2)$		
	$\text{Hsafe2: safe}_\varphi(e_2, \sigma_2)$		
<hr/>			
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$		
<hr/>			
Spatial hyp	$\text{Ho: S}_o(\sigma_1)$		
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Goal	$\text{wp } e_1 \{v. \lceil \varphi(v) \rceil\}$		

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[e'_1], \sigma_1)$	unfold prim-step	$e_1 = K[e'_1] \wedge e_2 = K[e'_2]$
	$\text{Hbase: } (e'_1, \sigma_1) \rightarrow_{\text{base}} (e'_2, \sigma_2)$		
	$\text{Hsafe2: safe}_\varphi(K[e'_2], \sigma_2)$		
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$		
Spatial hyp	$\text{Ho: S}_o(\sigma_1)$		
Goal	$\text{wp } K[e'_1] \{v. \lceil \varphi(v) \rceil\}$		

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp $\text{Hsafe1}: \text{safe}_\varphi(K[e'_1], \sigma_1)$

$\text{Hbase}: (e'_1, \sigma_1) \rightarrow_{\text{base}} (e'_2, \sigma_2)$

$\text{Hsafe2}: \text{safe}_\varphi(K[e'_2], \sigma_2)$

Persistent hyp $\text{IH}: \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp $\text{Ho}: \text{S}_o(\sigma_1)$

Goal $\text{wp } K[e'_1] \{v. \lceil \varphi(v) \rceil\}$

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Spatial hyp $\text{Ho: S}_o(\sigma_1)$

Goal $\text{wp } K[e'_1] \{v. \lceil \varphi(v) \rceil\}$

by WPBIND

*

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[e'_1], \sigma_1)$

Hbase: $(e'_1, \sigma_1) \rightarrow_{\text{base}} (e'_2, \sigma_2)$

Hsafe2: $\text{safe}_\varphi(K[e'_2], \sigma_2)$

Persistent hyp

IH: $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp

H_o: $\mathcal{S}_o(\sigma_1)$

Goal

$\text{wp } e'_1 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$ by WPBIND

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp $\text{Hsafe1: safe}_\varphi(K[e'_1], \sigma_1)$

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Spatial hyp $\text{Ho: S}_o(\sigma_1)$

Goal $\text{wp } e'_1 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp Hsafe1: $\text{safe}_\varphi(K[e'_1], \sigma_1)$ case analysis

Hbase: $(e'_1, \sigma_1) \rightarrow_{\text{base}} (e'_2, \sigma_2)$

Hsafe2: $\text{safe}_\varphi(K[e'_2], \sigma_2)$

Persistent hyp IH: $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp Ho: $\mathcal{S}_o(\sigma_1)$

Goal $\text{wp } e'_1 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$)

Pure hyp $\text{Hsafe1: safe}_\varphi(K[e'_1], \sigma_1)$ case analysis

$\text{Hbase: } (e'_1, \sigma_1) \rightarrow_{\text{base}} (e'_2, \sigma_2)$

$\text{Hsafe2: safe}_\varphi(K[e'_2], \sigma_2)$

Persistent hyp $\square e'_1 \rightarrow_{\text{pure}} e'_2, \sigma_1 \vdash \text{safe}_\varphi(e, \sigma) \rightarrow_{\text{base}} S_0(\sigma) \rightarrow_{\text{wp}} e \{v. \text{safe}_\varphi(v)\}$ \square

Spatial hyp $\square \sigma_1(l) = v \wedge (!l, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

$\square \sigma_1(l) = v \wedge (l \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((), \sigma_1[l \leftarrow w])$

Goal $\square l_0 \notin \sigma_1 \wedge (\text{ref } v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$ *

If e'_1 takes a pure step

Pure hyp $\text{Hsafe1: safe}_\varphi(K[e'_1], \sigma_1)$

$\text{Hpure: } e'_1 \rightarrow_{\text{pure}} e'_2$

$\text{Hsafe2: safe}_\varphi(K[e'_2], \sigma_1)$

Persistent hyp $\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp $\text{H}_o: \text{S}_o(\sigma_1)$

Goal $\text{wp } e'_1 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

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Spatial hyp $\text{Ho: S}_o(\sigma_1)$

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by **WPPURE**

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Spatial hyp $\text{H}_o: \text{S}_o(\sigma_1)$

Goal $\triangleright \text{wp } e'_2 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

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$H_{\text{pure}}: e'_1 \rightarrow_{\text{pure}} e'_2$

$H_{\text{safe2}}: \text{safe}_\varphi(K[e'_2], \sigma_1)$

Persistent hyp $IH: \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap S_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

strip later

Spatial hyp $H_o: S_o(\sigma_1)$

Goal $\triangleright \text{wp } e'_2 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

If e'_1 takes a pure step

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Spatial hyp $\text{Ho: S}_o(\sigma_1)$

Goal

$\text{wp } e'_2 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

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Spatial hyp $\text{Ho: S}_o(\sigma_1)$

Goal $\text{wp } e'_2 \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

by WPBINDINV

If e'_1 takes a pure step

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If e'_1 takes a pure step

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$H_{\text{pure}}: e'_1 \rightarrow_{\text{pure}} e'_2$

$H_{\text{safe2}}: \text{safe}_\varphi(K[e'_2], \sigma_1)$

Persistent hyp $IH: (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap S_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$ apply

Spatial hyp $H_o: S_o(\sigma_1)$

Goal $\text{wp } K[e'_2] \{v. \lceil \varphi(v) \rceil\}$ *

If e'_1 takes a pure step

Pure hyp $H_{\text{safe1}}: \text{safe}_\varphi(K[e'_1], \sigma_1)$

$H_{\text{pure}}: e'_1 \rightarrow_{\text{pure}} e'_2$

$H_{\text{safe2}}: \text{safe}_\varphi(K[e'_2], \sigma_1)$

Persistent hyp $IH: (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap S_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$ apply

Spatial hyp $H_o: S_o(\sigma_1)$

Goals $\lceil \text{safe}_\varphi(K[e'_2], \sigma_1) \rceil$ $S_o(\sigma_1)$ *

If e'_1 takes a pure step

Pure hyp $\text{Hsafe1: safe}_\varphi(K[e'_1], \sigma_1)$

$\text{Hpure: } e'_1 \rightarrow_{\text{pure}} e'_2$

$\text{Hsafe2: safe}_\varphi(K[e'_2], \sigma_1)$ **frame**

Persistent hyp $\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp $\text{Ho: S}_o(\sigma_1)$ **frame**

Goals $\lceil \text{safe}_\varphi(K[e'_2], \sigma_1) \rceil$ $\text{S}_o(\sigma_1)$ *

If e'_1 takes a pure step

Pure hyp Hsafe1: $\text{safe}_\varphi(K[e'_1], \sigma_1)$

✓ $e'_1 \rightarrow_{\text{pure}} e'_2$

□ $\sigma_1(l) = v \wedge (!l, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Persistent hyp □ $\sigma_1(l) = v \wedge (l \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((l), \sigma_1[l \leftarrow w])$

Spatial hyp □ $l_0 \notin \sigma_1 \wedge (\text{ref } v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$

Goal



*

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathcal{S}_o(\sigma_1)$	□
Goal	$\text{wp } !\ell \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$	

If $e'_1 = !\ell$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[!\ell], \sigma_1)$

Hlookup: $\sigma_1(\ell) = v$

Hbase: $(!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Hsafe2: $\text{safe}_\varphi(K[v], \sigma_1)$

Persistent hyp

IH: $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil * \mathcal{S}_o(\sigma) * \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp

H_o: $\mathcal{S}_o(\sigma_1)$

destruct

$\mathcal{S}_o(\sigma_1) \triangleq \bigstar_{(\ell \leftarrow v) \in \sigma_1} (\ell \mapsto v)$

Goal

$\text{wp } !\ell \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil * \mathcal{S}_o(\sigma) * \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{H0: } \mathcal{S}_o(\sigma_1 \setminus \{\ell\})$	$\text{destruct } \mathcal{S}_o(\sigma_1) \triangleq \bigstar_{(\ell \leftarrow v) \in \sigma_1} (\ell \mapsto v)$
	$\text{H1: } \ell \mapsto v$	
Goal	$\text{wp } !\ell \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$	

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
<hr/>		
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil * \mathbb{S}_o(\sigma) * \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
<hr/>		
Spatial hyp	$\text{Ho: } \mathbb{S}_o(\sigma_1 \setminus \{\ell\})$	
	$\text{Hl: } \ell \mapsto v$	
<hr/>		
Goal	$\text{wp } !\ell \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$	

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathbb{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathbb{S}_o(\sigma_1 \setminus \{\ell\})$	
	$\text{H1: } \ell \mapsto v$	
Goal	$\text{wp } !\ell \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$	
	by WPLOAD	

If $e'_1 = !\ell$

Pure hyp Hsafe1: $\text{safe}_\varphi(K[!\ell], \sigma_1)$ Hlookup: $\sigma_1(\ell) = v$

Hbase: $(!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Hsafe2: $\text{safe}_\varphi(K[v], \sigma_1)$

Persistent hyp IH: $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp Ho: $\mathcal{S}_o(\sigma_1 \setminus \{\ell\})$

Goal $\triangleright (\ell \mapsto v \multimap \text{wp } K[v] \{v. \lceil \varphi(v) \rceil\})$

by WPLOAD

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \text{S}_o(\sigma_1 \setminus \{\ell\})$	
Goal	$\triangleright (\ell \mapsto v \multimap \text{wp } K[v] \{v. \lceil \varphi(v) \rceil\})$	

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \text{S}_o(\sigma_1 \setminus \{\ell\})$	
Goal	$\triangleright (\ell \mapsto v \multimap \text{wp } K[v] \{v. \lceil \varphi(v) \rceil\})$	

strip later and intro

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathbb{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathbb{S}_o(\sigma_1 \setminus \{\ell\})$	
	$\text{Hl: } \ell \mapsto v$	
Goal	$\text{wp } K[v] \{v. \lceil \varphi(v) \rceil\}$	

strip later and intro

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil * \mathbb{S}_o(\sigma) * \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathbb{S}_o(\sigma_1 \setminus \{\ell\})$	
	$\text{Hl: } \ell \mapsto v$	
Goal	$\text{wp } K[v] \{v. \lceil \varphi(v) \rceil\}$	

If $e'_1 = !\ell$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[!\ell], \sigma_1)$

Hlookup: $\sigma_1(\ell) = v$

Hbase: $(!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Hsafe2: $\text{safe}_\varphi(K[v], \sigma_1)$

Persistent hyp

IH: $(\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_\circ(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp

H₀: $\mathcal{S}_\circ(\sigma_1 \setminus \{\ell\})$

combine

$\mathcal{S}_\circ(\sigma_1) \triangleq \bigstar_{(\ell \leftarrow v) \in \sigma_1} (\ell \mapsto v)$

H₁: $\ell \mapsto v$

Goal

$\text{wp } K[v] \{v. \lceil \varphi(v) \rceil\}$

If $e'_1 = !\ell$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[!\ell], \sigma_1)$

Hlookup: $\sigma_1(\ell) = v$

Hbase: $(!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Hsafe2: $\text{safe}_\varphi(K[v], \sigma_1)$

Persistent hyp

IH: $(\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_\circ(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp

$\text{Ho}: \mathcal{S}_\circ(\sigma_1)$

combine

$\mathcal{S}_\circ(\sigma_1) \triangleq \bigstar_{(l \leftarrow v) \in \sigma_1} (l \mapsto v)$

Goal

$\text{wp } K[v] \{v. \lceil \varphi(v) \rceil\}$

If $e'_1 = !\ell$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[!\ell], \sigma_1)$	$\text{Hlookup: } \sigma_1(\ell) = v$
	$\text{Hbase: } (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$	
	$\text{Hsafe2: safe}_\varphi(K[v], \sigma_1)$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \text{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \text{S}_o(\sigma_1)$	
Goal	$\text{wp } K[v] \{v. \lceil \varphi(v) \rceil\}$	

If $e'_1 = !\ell$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[!\ell], \sigma_1)$

Hlookup: $\sigma_1(\ell) = v$

Hbase: $(!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Hsafe2: $\text{safe}_\varphi(K[v], \sigma_1)$ frame

Persistent hyp

IH: $(\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil * S_o(\sigma) * \text{wp } e \{v. \lceil \varphi(v) \rceil\})$ apply

Spatial hyp

H_o: $S_o(\sigma_1)$ frame

Goal

$\text{wp } K[v] \{v. \lceil \varphi(v) \rceil\}$

*

If $e'_1 = !\ell$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[!\ell], \sigma_1)$

Hlookup: $\sigma_1(\ell) = v$

Hbase: $(!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Hsafe2: $\text{safe}_\varphi(K[v], \sigma_1)$

Persistent hyp

IH: $(\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil * S_o(\sigma) * \text{wp } e \{v. \lceil \varphi(v) \rceil\})$ apply

Spatial hyp

H_o: $S_o(\sigma_1)$ framed

Goal



*

If $e'_1 = !l$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[!l], \sigma_1)$

Hlookup: $\sigma_1(l) = v$

✓ $e'_1 \rightarrow_{\text{pure}} e'_2$

✓ $\sigma_1(l) = v \wedge (!l, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Persistent

□ $\sigma_1(l) = v \wedge (l \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((l), \sigma_1[l \leftarrow w])$

Spatial hyp

□ $l_0 \notin \sigma_1 \wedge (\text{ref } v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$

Goal



*

If $e'_1 = \ell \leftarrow w$

Pure hyp $\text{Hsafe1: safe}_\varphi(K[\ell \leftarrow w], \sigma_1)$ $\text{Hlookup: } \sigma_1(\ell) = v$

$\text{Hbase: } (\ell \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((), \sigma_1[\ell \leftarrow w])$

$\text{Hsafe2: safe}_\varphi(K[()], \sigma_1[\ell \leftarrow w])$

Persistent hyp $\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp $\text{Ho: } \mathcal{S}_o(\sigma_1)$

Goal $\text{wp } \ell \leftarrow w \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

If $e'_1 = \ell \leftarrow w$

Pure hyp Hsafe1: $\text{safe}_\varphi(K[\ell \leftarrow w], \sigma_1)$ Hlookup: $\sigma_1(\ell) = v$

Hbase: $(\ell \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((), \sigma_1[\ell \leftarrow w])$

Hsafe2: $\text{safe}_\varphi(K[()], \sigma_1[\ell \leftarrow w])$

Persistent hyp IH: $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp Ho: $\mathcal{S}_o(\sigma_1)$

Goal $\text{wp } \ell \leftarrow w \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

by WPSTORE

If $e'_1 = \ell \leftarrow w$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[\ell \leftarrow w], \sigma_1)$

Hlookup: $\sigma_1(\ell) = v$

✓ $e'_1 \rightarrow_{\text{pure}} e'_2$

✓ $\sigma_1(\ell) = v \wedge (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$

Persistent

✓ $\sigma_1(\ell) = v \wedge (\ell \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((\ell), \sigma_1[\ell \leftarrow w])$

Spatial hyp

□ $\ell_0 \notin \sigma_1 \wedge (\text{ref } v, \sigma_1) \rightarrow_{\text{base}} (\ell_0, \sigma_1[\ell_0 \leftarrow v])$

Goal

$\text{wp } \ell \leftarrow w \{w. \text{wp } K[w] \{v. \lceil \varphi(v) \rceil\}\}$

by WPSTORE

If $e'_1 = \mathbf{ref} v$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$ $\text{Hbase: } (\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (\ell_0, \sigma_1[\ell_0 \leftarrow v])$ $\text{Hsafe2: safe}_\varphi(K[\ell_0], \sigma_1[\ell_0 \leftarrow v])$
Persistent hyp	$\text{IH: } \triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$
Spatial hyp	$\text{H}_o: \mathcal{S}_o(\sigma_1)$
Goal	$\mathbf{wp} \mathbf{ref} v \{w. \mathbf{wp} K[w] \{v. \lceil \varphi(v) \rceil\}\}$

If $e'_1 = \mathbf{ref} v$

Pure hyp

Hsafe1: $\mathbf{safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$

Hbase: $(\mathbf{ref} v, \sigma_1) \rightarrow_{\mathbf{base}} (\ell_0, \sigma_1[\ell_0 \leftarrow v])$

Hsafe2: $\mathbf{safe}_\varphi(K[\ell_0], \sigma_1[\ell_0 \leftarrow v])$

Persistent hyp

IH: $\triangleright (\forall e, \sigma. \lceil \mathbf{safe}_\varphi(e, \sigma) \rceil \multimap \mathbf{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp

H_o: $\mathbf{S}_o(\sigma_1)$

Goal

$\mathbf{wp} \mathbf{ref} v \{w. \mathbf{wp} K[w] \{v. \lceil \varphi(v) \rceil\}\}$

by WPALLOC

If $e'_1 = \mathbf{ref} v$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$	$\text{Hfresh: } l \notin \sigma_1$
	$\text{Hbase: } (\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$	
	$\text{Hsafe2: safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathcal{S}_o(\sigma_1[l \leftarrow v])$	
Goal	$\mathbf{wp} K[l] \{v. \lceil \varphi(v) \rceil\}$	

by WPALLOC

If $e'_1 = \mathbf{ref} v$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$	$\text{Hfresh: } l \notin \sigma_1$
	$\text{Hbase: } (\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$	
	$\text{Hsafe2: safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathcal{S}_o(\sigma_1[l \leftarrow v])$	
Goal	$\mathbf{wp} K[l] \{v. \lceil \varphi(v) \rceil\}$	

If $e'_1 = \mathbf{ref} v$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$	$\text{Hfresh: } l \notin \sigma_1$
	$\text{Hbase: } (\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$	$l_0 \text{ and } l \text{ are unrelated}$
	$\text{Hsafe2: safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathcal{S}_o(\sigma_1[l \leftarrow v])$	
Goal	$\mathbf{wp} K[l] \{v. \lceil \varphi(v) \rceil\}$	

If $e'_1 = \mathbf{ref} v$

Pure hyp	$\text{Hsafe1: safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$	$\text{Hfresh: } l \notin \sigma_1$
	$\text{Hbase: } (\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$	$l_0 \text{ and } l \text{ are unrelated}$
	$\text{Hsafe2: safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$	
Persistent hyp	$\text{IH: } (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap \mathcal{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$	
Spatial hyp	$\text{Ho: } \mathcal{S}_o(\sigma_1[l \leftarrow v])$	
Goal	$\mathbf{wp} K[l] \{v. \lceil \varphi(v) \rceil\}$	

If $e'_1 = \mathbf{ref} v$

Pure hyp Hsafe1: $\mathbf{safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$ Hfresh: $l \notin \sigma_1$

Hbase: $(\mathbf{ref} v, \sigma_1) \rightarrow_{\mathbf{base}} (l_0, \sigma_1[l_0 \leftarrow v])$

Hsafe2: $\mathbf{safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$

Persistent hyp IH: $(\forall e, \sigma. \lceil \mathbf{safe}_\varphi(e, \sigma) \rceil \multimap \mathbf{S}_o(\sigma) \multimap \mathbf{wp} e \{v. \lceil \varphi(v) \rceil\})$

Spatial hyp

Goal

$\lceil \mathbf{safe}_\varphi(K[l], \sigma_1[l \leftarrow v]) \rceil$

If $e'_1 = \mathbf{ref} v$

Pure hyp Hsafe1: $\mathit{safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$ Hfresh: $l \notin \sigma_1$

pure intro

Hbase: $(\mathbf{ref} v, \sigma_1) \rightarrow_{\mathit{base}} (l_0, \sigma_1[l_0 \leftarrow v])$

Hsafe2: $\mathit{safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$

Persistent hyp IH: $(\forall e, \sigma. \lceil \mathit{safe}_\varphi(e, \sigma) \rceil \multimap \mathit{S}_o(\sigma) \multimap \mathit{wp} e \{v. \lceil \varphi(v) \rceil\})$

□

Spatial hyp

Goal

$\lceil \mathit{safe}_\varphi(K[l], \sigma_1[l \leftarrow v]) \rceil$

*

If $e'_1 = \mathbf{ref} v$

Pure hyp

Hsafe1: $\text{safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$

Hfresh: $l \notin \sigma_1$

Hbase: $(\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (l_0, \sigma_1[l_0 \leftarrow v])$

Hsafe2: $\text{safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$

$\text{safe}_\varphi(K[l], \sigma_1[l \leftarrow v])$

If $e'_1 = \mathbf{ref} v$

Pure hyp $\mathbf{Hsafe1}: \mathbf{safe}_\varphi(K[\mathbf{ref} v], \sigma_1)$ $\mathbf{Hfresh}: l \notin \sigma_1$

$\mathbf{Hbase}: (\mathbf{ref} v, \sigma_1) \rightarrow_{\mathbf{base}} (l_0, \sigma_1[l_0 \leftarrow v])$

$\mathbf{Hsafe2}: \mathbf{safe}_\varphi(K[l_0], \sigma_1[l_0 \leftarrow v])$

$\mathbf{safe}_\varphi(K[l], \sigma_1[l \leftarrow v])$

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$(\mathbf{ref} v, \sigma_1) \rightarrow_{\mathbf{base}} (l, \sigma_1[l \leftarrow v])$

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$(\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (l, \sigma_1[l \leftarrow v])$ ✓

(because **ref** chooses l non-deterministically)

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$\text{safe}_\varphi(K[l], \sigma_1[l \leftarrow v])$

\uparrow

✓ $e'_1 \rightarrow_{\text{pure}} e'_2 \rightarrow_{\text{base}} (\ell, \sigma_1[l \leftarrow v])$ ✓

✓ $\sigma_1(l) = v \wedge (!\ell, \sigma_1) \rightarrow_{\text{base}} (v, \sigma_1)$ ℓ non-deterministically

✓ $\sigma_1(l) = v \wedge (\ell \leftarrow w, \sigma_1) \rightarrow_{\text{base}} ((), \sigma_1[l \leftarrow w])$

✓ $l_0 \notin \sigma_1 \wedge (\mathbf{ref} v, \sigma_1) \rightarrow_{\text{base}} (\ell_0, \sigma_1[l_0 \leftarrow v])$

Recap

- Löb induction $\triangleright (\forall e, \sigma. \lceil \text{safe}_\varphi(e, \sigma) \rceil \multimap S_o(\sigma) \multimap \text{wp } e \{v. \lceil \varphi(v) \rceil\})$
- If $\text{safe}_\varphi(v, \sigma)$ ($v \in \text{Val}$) $\implies \lceil \varphi(v) \rceil \vdash \text{wp } v \{v. \lceil \varphi(v) \rceil\}$ (by WPVALUE)
- If $\text{safe}_\varphi(e_1, \sigma_1)$ ($e_1 \notin \text{Val}$), then e_1 is reducible

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case analysis on the first step

Language-independent, expect for the blue part

For language-dependent part: update $S_o(\sigma)$ using primitive laws.

Concurrent Completeness

Sequential:

- Proof environment: $\text{safe}_\varphi(e, \sigma)$ and $S_o(\sigma)$
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$$I_{\text{compl}} \triangleq \boxed{\exists \vec{e}, \sigma. \lceil \text{safe-tp}_\varphi(\vec{e}, \sigma) \rceil * S_o(\sigma)}^{\mathcal{N}}$$
$$\forall i, e. I_{\text{compl}} \quad \vdash \text{wp } e \{v. \lceil i = 0 \Rightarrow \varphi(v) \rceil\}$$

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Generalization

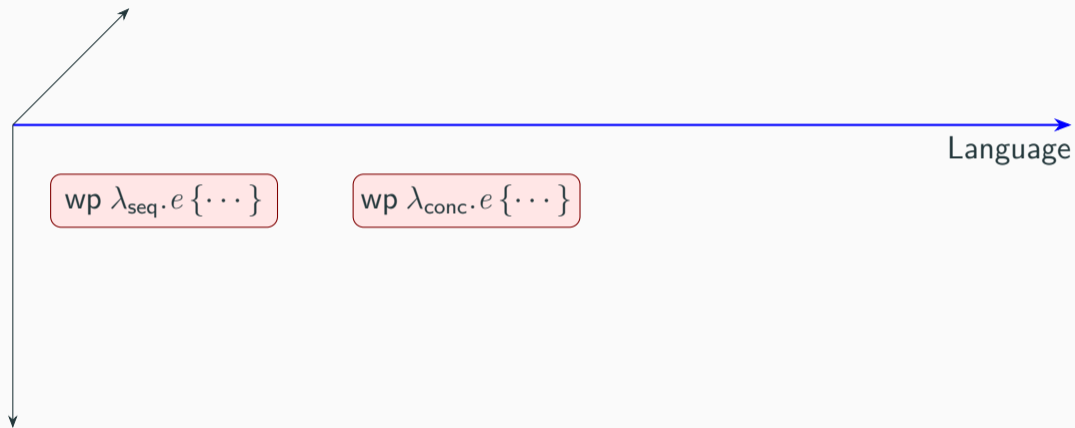
$\text{wp } \lambda_{\text{seq}} \cdot e \{ \dots \}$

$\text{wp } \lambda_{\text{conc}} \cdot e \{ \dots \}$

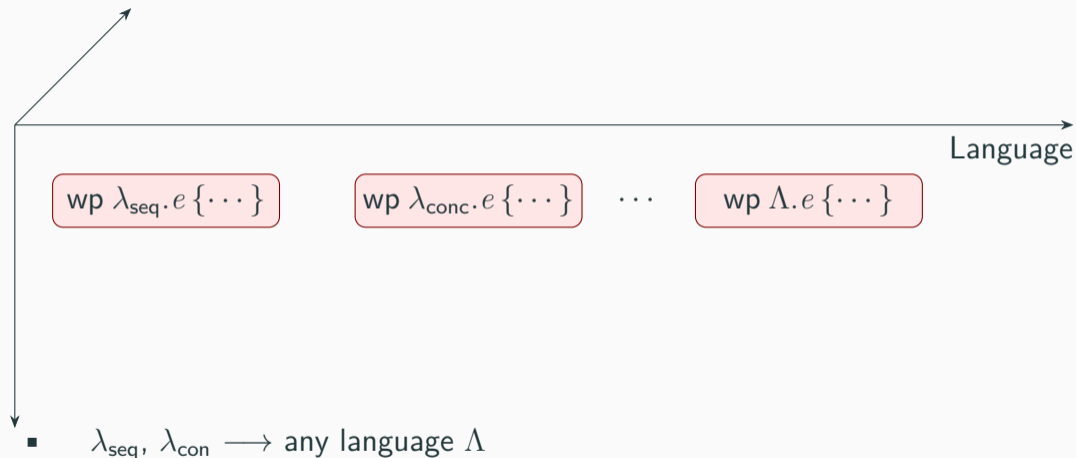
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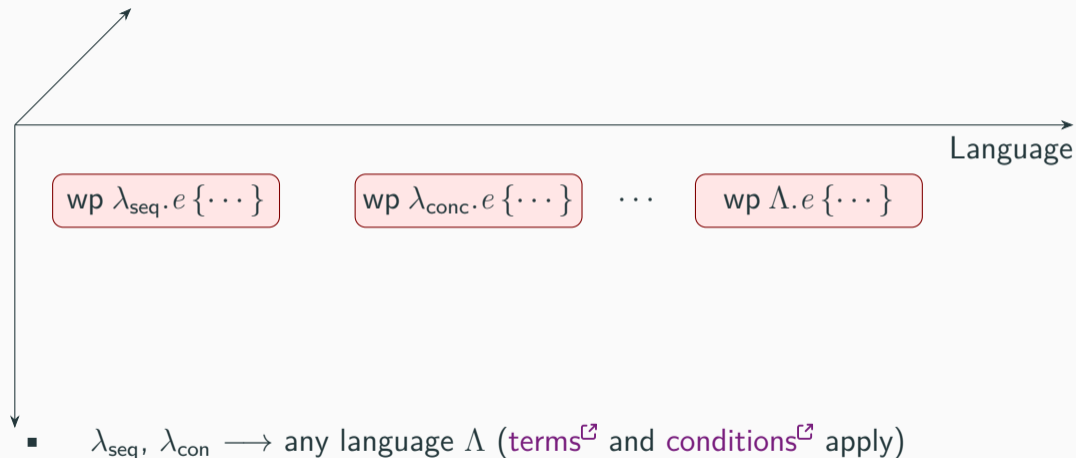
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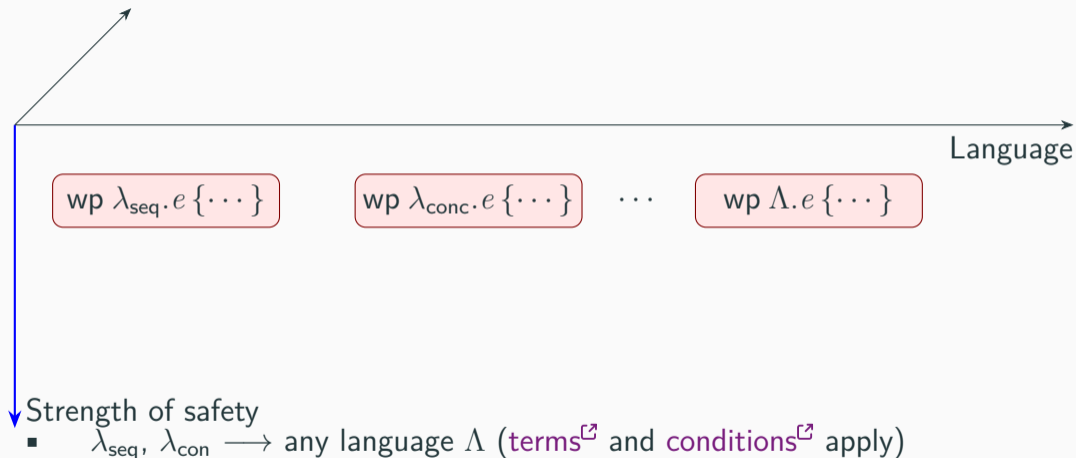
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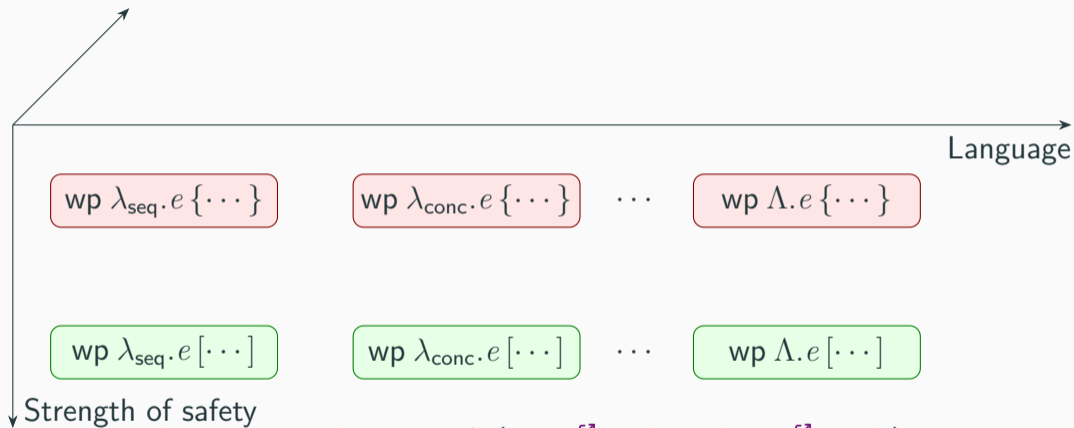
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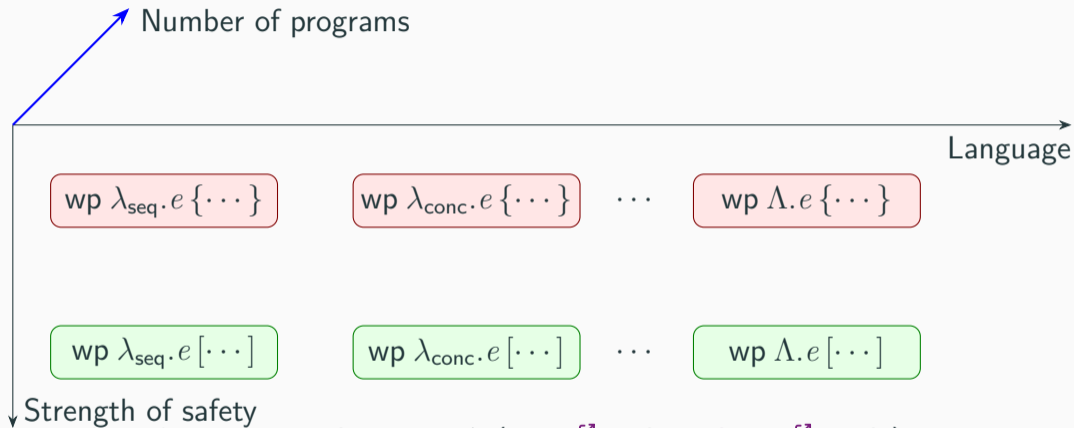


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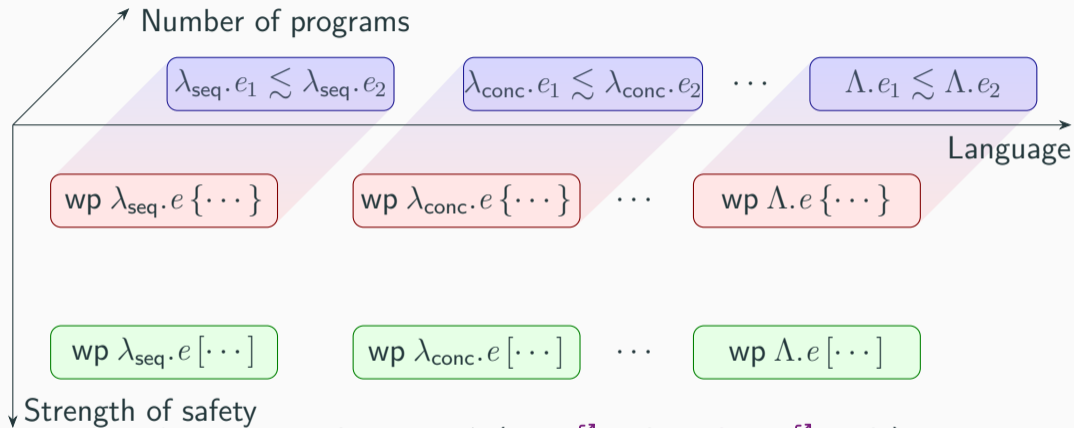
- $\lambda_{\text{seq}}, \lambda_{\text{con}} \longrightarrow$ any language Λ (terms[↗] and conditions[↗] apply)
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Condition 32 (Completeness).

$$\lceil \text{base-red}(e, \sigma) \rceil * \mathbb{S}_\circ(\sigma) * \bullet^Y \vec{e} * \lceil \text{safe-tp}_\varphi(\vec{e}, \sigma) \rceil * n \hookrightarrow^Y K[e] \vdash \models_{\mathcal{E}}$$

$$(\lceil \text{Atomic } e \rceil * \forall \Phi. \triangleright \text{WithStatePre}(\mathcal{E}, e, n, K, \sigma, \vec{e}, \Phi) * \text{wp}_{\mathcal{E}} e \{\Phi\}) \vee$$

$$(\mathbb{S}_\circ(\sigma) * \bullet^Y \vec{e} * \forall \Phi. \triangleright \text{WithoutStatePre}(\mathcal{E}, e, K, \Phi) * \text{wp}_{\top} e \{\Phi\})$$

where

$$\text{WithStatePre}(\mathcal{E}, e, n, K, \sigma, \vec{e}, \Phi) \triangleq \forall v', \sigma', \vec{e}_f. \lceil (e, \sigma) \rightarrow (v', \sigma', \vec{e}_f) \rceil * \mathbb{S}_\circ(\sigma') * \bullet^Y \vec{e} * n \hookrightarrow^Y K[e] * \models_{\mathcal{E}}$$

$$\Phi(v') * *_{e_f \in \vec{e}_f} \text{wp}_{\top} e_f \{v. \text{True}\}$$

$$\text{WithoutStatePre}(\mathcal{E}, e, K, \Phi) \triangleq \forall e', \vec{e}, \vec{e}_f. \text{Reduces}(e, K, e', \vec{e}) * \models_{\top}$$

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$$\exists \sigma'. \lceil (e, \sigma) \rightarrow^+ (e', \sigma', \vec{e}_f) \rceil * \mathbb{S}_\circ(\sigma') * \bullet^Y \vec{e} * n \hookrightarrow^Y K[e]$$

wp λ_{se} $\approx \Lambda.e_2$

Language

Completeness. If Condition 32 holds for a language Λ , then

$$\text{safe}_\varphi(e, \emptyset) \implies \vdash \text{wp } e \{v. \lceil \varphi(v) \rceil\}$$

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- λ_{Rust} : handle non-atomic memory access
- Eris: redo the completeness theorem following the same pattern

The missing rules in HeapLang

```
unique  $\triangleq$  let  $\ell_1 = \mathbf{ref}()$  in  
    free  $\ell_1$ ;  
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```

The missing rules in HeapLang

$unique \triangleq$ **let** $l_1 = \mathbf{ref}()$ **in**
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$$\frac{l \notin \text{dom}(\sigma)}{(\mathbf{ref} v, \sigma) \rightarrow_{\text{base}} (l, \sigma[l \leftarrow v], \epsilon)}$$

Safe because locations are never reused!

$$\frac{\sigma(l) = v}{(\mathbf{free} l, \sigma) \rightarrow_{\text{base}} ((), \sigma[l \leftarrow \perp], \sigma, \epsilon)}$$

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Another Example

resolve \triangleq **resolve** (LitProph 1) **to** ()

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Completeness.

For a HeapLang expression e , $\text{safe}_\varphi(e, \emptyset) \implies \vdash \text{wp } e \{v. \lceil \varphi(v) \rceil\}$

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Insight: What Iris features are needed?

- Löb induction
- First-order ghost state, ghost maps
- Invariants of timeless assertions
- `WPVALUE`, `WPBIND`, `WPPURE`, `FUPDWP`, and `WPATOMIC`

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- **Locality of the semantics**
 - Allocation must non-deterministically pick a location
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We don't need

- Higher-order ghost state
- Impredicative invariants
- Later credits

Concurrent Library Completeness

Zichen Zhang

Simon Gregersen

Joseph Tassarotti

Logic & Truth for Concurrent Data Structures

Logic

Truth

Logic & Truth for Concurrent Data Structures

Logic

$\text{lin}_\mu(\text{op})$

Linearizable *w.r.t.*
sequential behavior μ

Truth

Logic & Truth for Concurrent Data Structures

$I_\mu \vdash \langle P_\mu \rangle \text{op}(obj, x) \langle y. Q_\mu(y) \rangle$

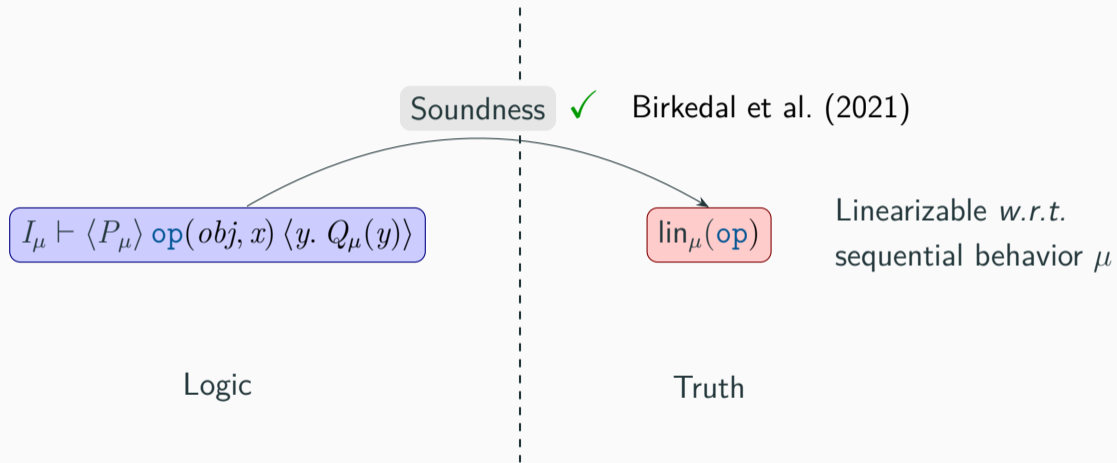
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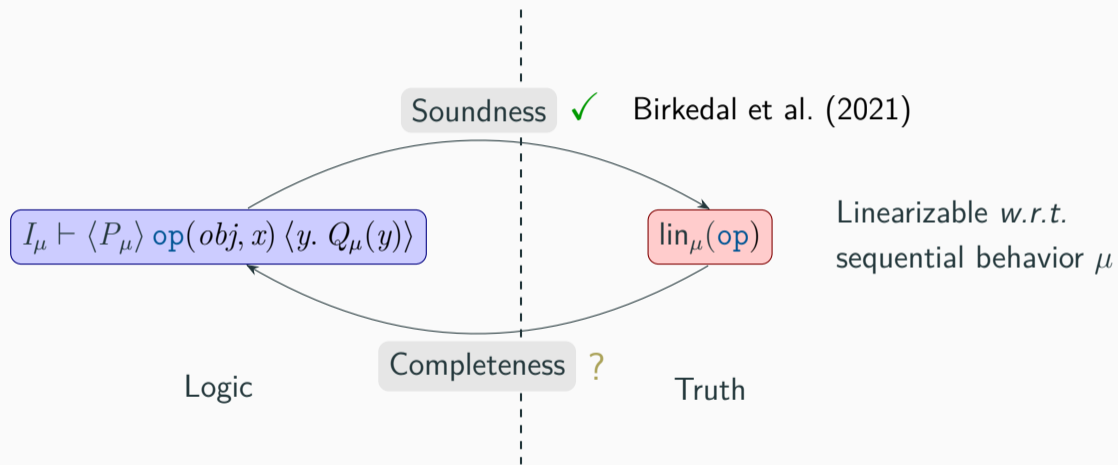
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Logic & Truth for Concurrent Data Structures

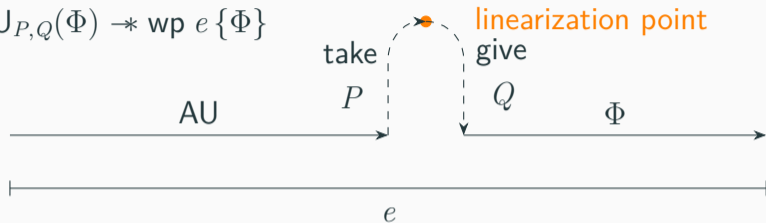


Logic & Truth for Concurrent Data Structures



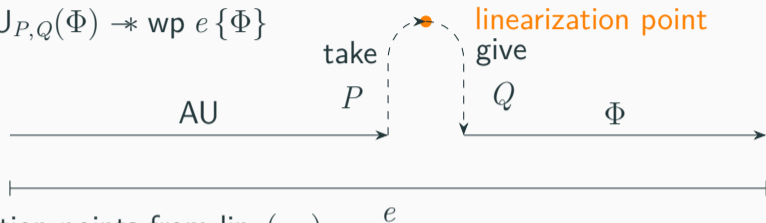
Key Ideas

$$\langle P \rangle e \langle Q \rangle \triangleq \forall \Phi. \text{AU}_{P,Q}(\Phi) \rightarrow * \text{wp } e \{ \Phi \}$$



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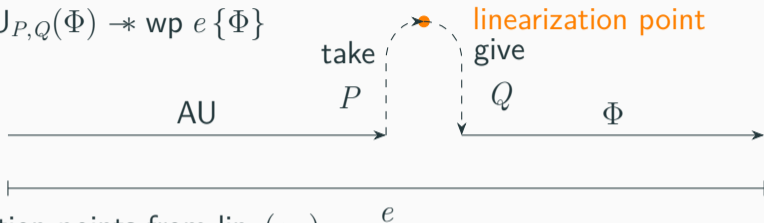
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- Extract linearization points from $\text{lin}_\mu(\text{op})$
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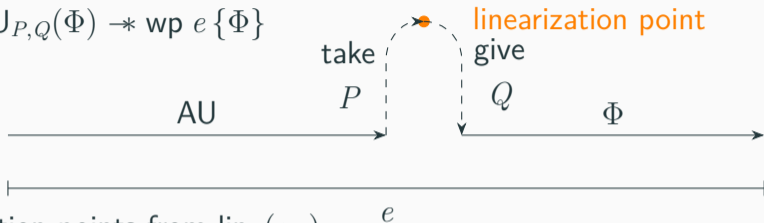
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- Future dependency:

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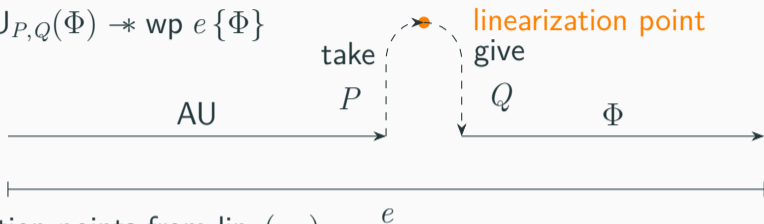
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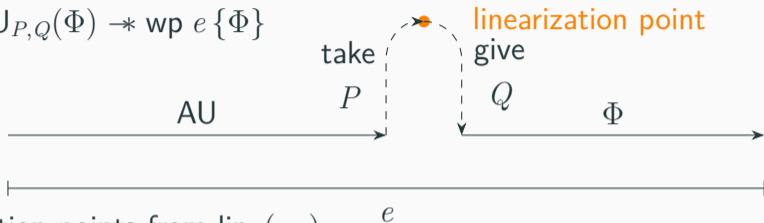
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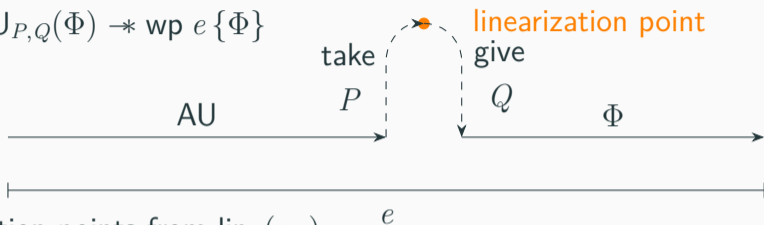
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- Future dependency: typed prophecy variables on undecidable types
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- Non structural linearization points: helping
needs impredicative invariants and later credits

Key Ideas

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Completeness.

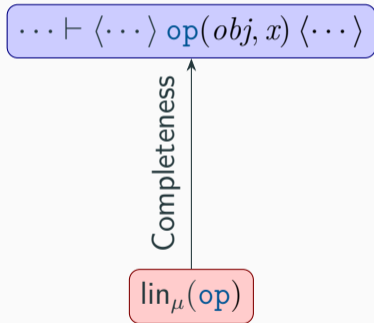
For a library op , $\text{lin}_\mu(\text{op}) \implies \forall \text{obj}, x. I_\mu \vdash \langle P_\mu \rangle \text{op}(\text{obj}, x) \langle y. Q_\mu(y) \rangle$

Challenges & Strategies

- Future dependency: typed prophecy variables on undecidable types
needs excluded middle law
- Non structural linearization points: helping
needs impredicative invariants and later credits

Why completeness?

- Theoretically interesting



Why completeness?

- Theoretically interesting
- Embedding external techniques

$\dots \vdash \langle \dots \rangle \text{op}(obj, x) \langle \dots \rangle$

Completeness



$\text{lin}_{\mu}(\text{op})$

Why completeness?

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$\dots \vdash \langle \dots \rangle \text{op}(obj, x) \langle \dots \rangle$

Completeness

$\text{lin}_\mu(\text{op})$

Aspect-oriented^[Henzinger et al.(2013)]

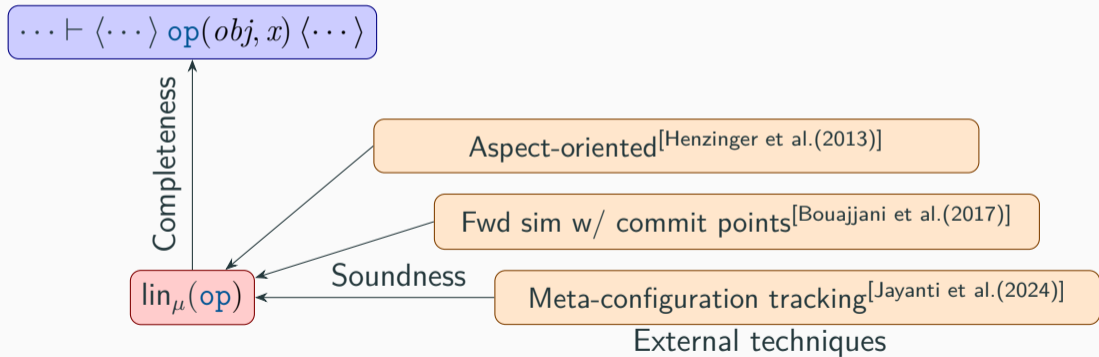
Fwd sim w/ commit points^[Bouajjani et al.(2017)]

Meta-configuration tracking^[Jayanti et al.(2024)]

External techniques

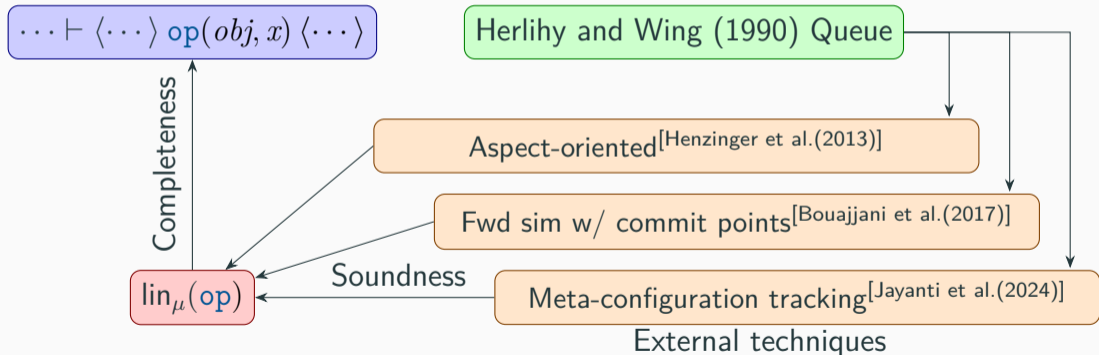
Why completeness?

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Conclusion

Part 1.

- Language-independent whole program completeness theorems
- Case studies: HeapLang, λ_{Rust} , time credits, and Eris

Conclusion

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- Language-independent whole program completeness theorems
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Part 2.

- Completeness of logically atomic triples
- Ported external techniques: aspect-oriented proofs, forward simulation with commit points, and meta-configuration tracking

Conclusion

Part 1. (To appear at ICFP 2026)

- Language-independent whole program completeness theorems
- Case studies: HeapLang, λ_{Rust} , time credits, and Eris




Please try it out!

(and/or join our discussion session)

<https://plf.inf.ethz.ch/research/icfp26-completeness.html>

Part 2.


- Completeness of logically atomic triples
- Ported external techniques: aspect-oriented proofs, forward simulation with commit points, and meta-configuration tracking

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